

Xeledon: Efficient enumeration of all steady state cycles in metabolic networks

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Genome-scale metabolic models summarize the knowledge about an organism and thus allow predictions about its features and capabilities. Most widely used are constraint-based models which rely on a set of chemical reactions and optionally include bounds for the flux through them. Validation and quality assurance are of great importance if the model of an organism is intended to be used for making predictions about the organisms' features or to guide expensive experiments. While the best way is validation against biological data, it is not available in all cases. Here, we address a common error occurring during a metabolic reconstruction: the inclusion of thermodynamically implausible reactions. Such reactions often lead to internal cycles which violate the second law of thermodynamics.^{4,6} We exploit this feature by enumerating all cycles in a metabolic network and so ease the search for thermodynamically implausible reactions. Furthermore, a cycle may be a feature of the organism's metabolism. Specifically, it has been shown that a futile cycle can regulate the equilibrium of a metabolite pool.³ From a regulatory point of view, cycles can be used to produce a feedback signal.¹ In the context of extreme pathways the cycles studied here have been classified as *Type III pathways*.⁵

In this work, we propose a fast algorithm which finds all steady state cycles in a metabolic network (algorithm 1). The key idea is to remove as many reactions as possible before carrying out the computationally expensive convex analysis. Since we are interested in internal fluxes only, all boundary reactions are removed first (lines 1 - 4). This creates blocked reactions which cannot carry any flux without violating the steady-state condition. Obviously, these reactions cannot participate in any cycle and thus are removed in line 5. Some reactions may have a bound which forces them to carry a non-zero flux. This may prevent the detection of unlimited reactions if they either cannot be active at all or their corresponding cycle cannot operate at the same time as some other cycle in the network. Thus, all flux bounds are relaxed in lines 7 and 8 so that they can be zero. All reactions which can still carry a flux at this point must be part of a cycle. Hence, all other reactions are identified and removed in line 9 using fast FVA.² Finally, the flux cone described by the remaining reactions is constructed in line 14 and all internal extreme pathways are enumerated. An example of an identified cycle caused by a PEP synthase running backwards is shown in figure 1.

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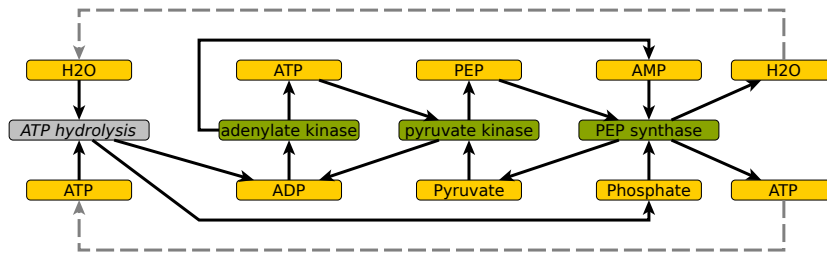


Figure 1: Example cycle which could regenerate ATP caused by a missing lower bound for the PEP synthase reaction

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Input : A set  $R$  of reactions
          Two vectors  $\vec{u}$  and  $\vec{l}$  with upper and lower flux
          bounds for each reaction in  $R$ 
          A set  $O$  of objective reactions
Output: A set  $C$  of sets of reactions which form a cycle

// remove all objective reactions
1 remove each reaction from  $R$  if reaction  $\in O$ 
2 foreach reaction in  $R$  do
3   | if reaction has no products or reaction has no educts
4   |   | then
5   |   |   | remove reaction from  $R$ 
6   |   |
7   |   | foreach reaction in  $R$  do
8   |   |   | if reaction is blocked then remove reaction from  $R$ 
9   |   |
10  |   | // allow each reaction to have zero flux
11  |   | Set each element  $e$  of  $\vec{u}$  to zero if  $e < 0$ 
12  |   | Set each element  $e$  of  $\vec{l}$  to zero if  $e > 0$ 
13  |   | // use FVA to remove reactions which cannot
14  |   |   | carry any flux
15  |   |   |  $\overrightarrow{\min flux}, \overrightarrow{\max flux} = \text{fastFVA}(R, \vec{u}, \vec{l})$ 
16  |   |   | foreach reaction in  $R$  do
17  |   |   |   | if  $\overrightarrow{\min flux}_{\text{reaction}} = 0$  and  $\overrightarrow{\max flux}_{\text{reaction}} = 0$  then
18  |   |   |   |   | remove reaction from  $R$ 
19  |   |   |
20  |   |   | // identify cycles using convex analysis
21  |   |   |  $C = \emptyset$ 
22  |   |   | foreach ray in  $\text{convexAnalysis}(R)$  do
23  |   |   |   | create a set  $C^r$  of reactions corresponding to the ray
24  |   |   |   | add  $C^r$  to  $C$ 

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Algorithm 1: xeledon algorithm

References

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